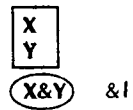


PRIMITIVE SENTENCE LOGIC RULES FOR NATURAL DEDUCTION

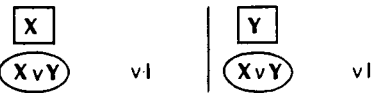
Conjunction Introduction



Conjunction Elimination



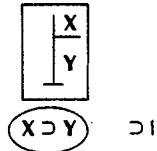
Disjunction Introduction



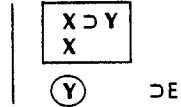
Disjunction Elimination



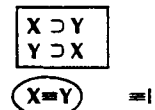
Conditional Introduction



Conditional Elimination



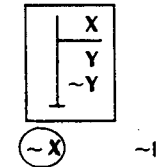
Biconditional Introduction



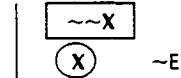
Biconditional Elimination



Negation Introduction

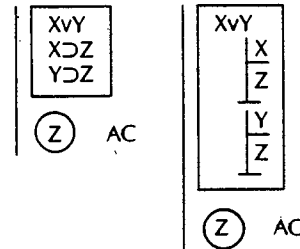


Negation Elimination

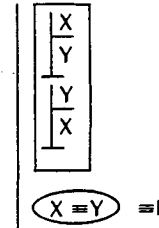


DERIVED SENTENCE LOGIC RULES FOR NATURAL DEDUCTION

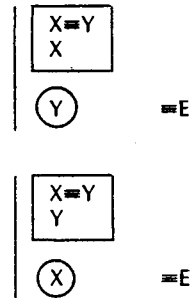
Argument by Cases



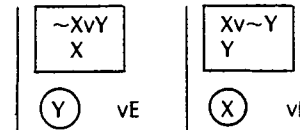
Biconditional Introduction



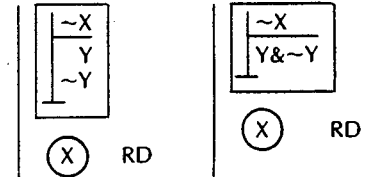
Biconditional Elimination



Disjunction Elimination



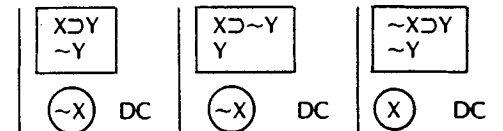
Reductio Ad Absurdum



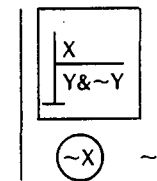
Weakening



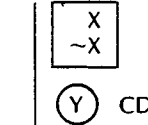
Denying the Consequent



Negation Introduction



Contradiction



De Morgan's Rules

- $\sim(X \vee Y)$ and $\sim X \wedge \sim Y$ are mutually derivable (DM)
- $\sim(X \wedge Y)$ and $\sim X \vee \sim Y$ are mutually derivable (DM)

Contraposition

- $X \supset Y$ and $\sim Y \supset \sim X$ are mutually derivable (CP)
- $\sim X \supset Y$ and $\sim Y \supset X$ are mutually derivable (CP)
- $X \supset \sim Y$ and $Y \supset \sim X$ are mutually derivable (CP)

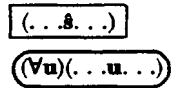
Conditional Rules

- $X \supset Y$ and $\sim X \vee Y$ are mutually derivable (C)
- $\sim(X \supset Y)$ and $X \wedge \sim Y$ are mutually derivable (C)

Reiteration: If a sentence occurs, either as a premise or as a conclusion in a derivation, that sentence may be copied (reiterated) in any of that derivation's LOWER sub-derivations, or lower down in the same derivation.

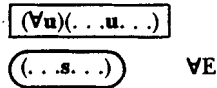
PRIMITIVE PREDICATE LOGIC RULES FOR NATURAL DEDUCTION

Universal Introduction

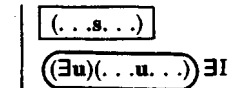


where \hat{s} occurs arbitrarily and $(\forall u)(\dots u \dots)$ is the universal generalization of $(\dots \hat{s} \dots)$

Universal Elimination

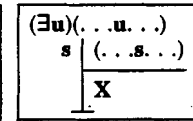


Existential Introduction



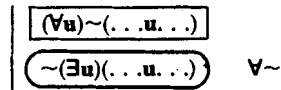
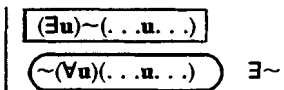
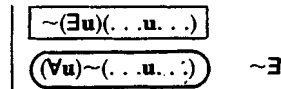
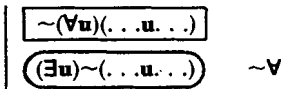
Where $(\exists u)(\dots u \dots)$ is an existential generalization of $(\dots s \dots)$

Existential Elimination

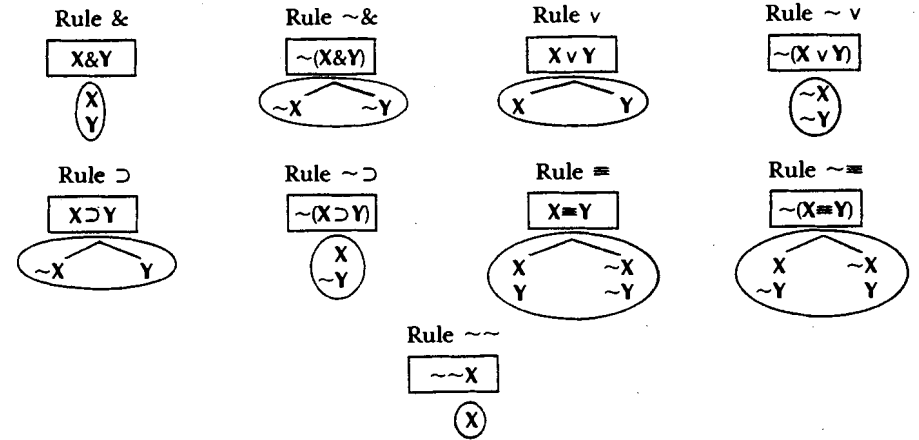


Where $(\dots s \dots)$ is a substitution instance of $(\exists u)(\dots u \dots)$ and s is isolated in the subderivation

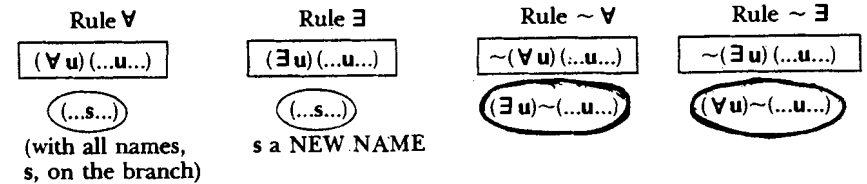
DERIVED RULES FOR NEGATED QUANTIFIERS



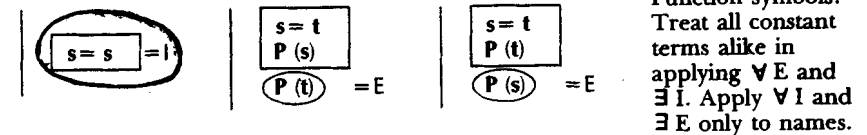
SENTENCE LOGIC TRUTH TREE RULES



PREDICATE LOGIC TRUTH TREE RULES



NATURAL DEDUCTION RULES FOR IDENTITY AND FUNCTION SYMBOLS



TRUTH TREE RULES FOR IDENTITY AND FUNCTION SYMBOLS

Rule \neq : Close any branch on which $s \neq s$ appears.
 Rule $=$: When $s = t$ appears on a branch, substitute s and t for each other wherever possible, without checking the resulting lines.

For function symbols: Instantiate all universally quantified sentences with all constant terms on the branch. In existentially quantified sentences use only one NEW NAME.